

Fast Pure Exploration via Frank-Wolfe

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Pure exploration on structured bandits

Stochastic Multi-Armed Bandit (MAB)

K arms (K prob. distribution ν_1, \dots, ν_K), the mean of ν_k is μ_k



ν_1



ν_2



ν_3



ν_4



ν_5

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In round t , an agent

1. pulls arm $A_t \in [K]$
2. receives the reward $X_{A_t}(t) \sim \nu_{A_t}$

Sequential sampling strategy: $A_t \in \mathcal{F}_t = \sigma[A_1, X_1, \dots, A_{t-1}, X_{t-1}]$

Pure exploration with fixed confidence

Goal: Identify a certain answer $i^*(\boldsymbol{\mu}) \in \mathcal{I}$

Example: Identify the best arm $i^*(\boldsymbol{\mu}) = \operatorname{argmax}_{k \in [K]} \mu_k$

A strategy consists of

- a sampling rule A_t (arm to explore)
- a stopping rule τ (time to stop)
- a \mathcal{F}_τ -measurable decision rule $\hat{i} \in \mathcal{I}$ (answer to return)

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We wish to minimize $\mathbb{E}_\mu[\tau]$ subject to $\mathbb{P}_\mu[\hat{i} \neq i^*(\boldsymbol{\mu})] < \delta$



“Side information” is encoded by the **structure**

Popular structures: Unstructured, Linear, Lipschitz, Dueling,
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Question 1. What is the sample complex gain achievable when exploiting the structure?

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Question 1. What is the sample complex gain achievable when exploiting the structure?

Question 2. Can we devise a computational efficient algorithm achieving the promised gains for all structures?

Lower bound [GK16]

For any **good** strategy,

$$\liminf_{\delta \rightarrow 0} \frac{\mathbb{E}_{\mu}[\tau]}{\log(\frac{1}{\delta})} \geq T^*(\mu),$$

where $T^*(\mu)^{-1} = \sup_{\omega \in \Sigma} \inf_{\lambda \in \text{Alt}(\mu)} \sum_{k=1}^K \omega_k d(\mu_k, \lambda_k)$

- Σ : $K - 1$ simplex
- $\text{Alt}(\mu) = \{\lambda \in \Lambda : i^*(\lambda) \neq i^*(\mu)\}$
- $d(\mu_k, \lambda_k)$: KL-divergent of arm-k reward distribution under λ and μ

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\Rightarrow An optimal algorithm has a sampling strategy described by

$$\omega^*(\mu) \in \operatorname{argmax}_{\omega \in \Sigma} F_{\mu}(\omega),$$

$$\text{where } F_{\mu}(\omega) = \inf_{\lambda \in \text{Alt}(\mu)} \sum_{k=1}^K \omega_k d(\mu_k, \lambda_k).$$



Generalized Likelihood Ratio Test (GLRT)

For each $k \in [K]$, $t \geq 1$, denote

- $N_k(t) = \sum_{s=1}^t \mathbb{1}\{A_s = k\}$,
- $\omega_k(t) = N_k(t)/t$,
- $\hat{\mu}_k(t) = \sum_{s=1}^t X_k(s) \mathbb{1}\{A_s = k\} / N_k(t)$ (when $N_k(t) > 0$),

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GLRT is the stopping rule s.t.

$\tau = \inf\{t \geq 1 : tF_{\hat{\mu}(t)}(\omega(t)) \geq \beta(t, \delta)\}$, where $\beta(t, \delta)$ satisfies:

$\forall t \geq 1, \left(tF_{\hat{\mu}(t)}(\omega(t)) \geq \beta(t, \delta) \right) \implies (\mathbb{P}_{\mu} [i^*(\hat{\mu}(t)) \neq i^*(\mu)] \leq \delta),$

$\exists c_1(\Lambda), c_2(\Lambda) > 0 : \forall t \geq c_1(\Lambda), \beta(t, \delta) \leq \log \left(\frac{c_2(\Lambda)t}{\delta} \right).$



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Many literatures[GK16, KK18, JP20, Mén19] provide such $\beta(t, \delta)$



To reach the optimality

Challenges for sampling rules:

- (i). μ is unknown initially
- (ii). **No** oracle for $\max_{\omega \in \Sigma} \inf_{\lambda \in \text{Alt}(\mu)} \sum_{k=1}^K \omega_k d(\mu_k, \lambda_k)$ in general

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For (ii)., Previous works either solve maxmini program for a single problem or converge to the saddle with overly conservative approach[Mén19, DKM19]



Frank-Wolfe based sampling (FWS)

Best challenger (BC) in unstructured bandits [GK16, Mén19]

Let $i^* = i^*(\boldsymbol{\mu})$ and $\text{Alt}(\boldsymbol{\mu}) = \cup_{j \neq i^*} \{\boldsymbol{\lambda} \in \Lambda : \lambda_j \geq \lambda_{i^*}\}$, then $F_{\boldsymbol{\mu}}(\boldsymbol{\omega}) = \min_{j \neq i^*} f_j(\boldsymbol{\omega}, \boldsymbol{\mu})$, where

$$\begin{aligned} f_j(\boldsymbol{\omega}, \boldsymbol{\mu}) &= \inf_{\lambda_j \geq \lambda_{i^*}} \sum_{k=1}^K \omega_k d(\mu_k, \lambda_k) \\ &= \omega_j d\left(\mu_j, \frac{\omega_{i^*} \mu_{i^*} + \omega_j \mu_j}{\omega_{i^*} + \omega_j}\right) + \omega_{i^*} d\left(\mu_{i^*}, \frac{\omega_{i^*} \mu_{i^*} + \omega_j \mu_j}{\omega_{i^*} + \omega_j}\right) \end{aligned}$$

Also, $\nabla_{\boldsymbol{\omega}} f_j(\boldsymbol{\omega}, \boldsymbol{\mu}) = d\left(\mu_j, \frac{\omega_{i^*} \mu_{i^*} + \omega_j \mu_j}{\omega_{i^*} + \omega_j}\right) \mathbf{e}_j + d\left(\mu_{i^*}, \frac{\omega_{i^*} \mu_{i^*} + \omega_j \mu_j}{\omega_{i^*} + \omega_j}\right) \mathbf{e}_{i^*}$

After pulling each arm once, BC repeatedly does:

1. Assign $C_t \leftarrow \operatorname{argmin}_{j \neq i^*}(\hat{\boldsymbol{\mu}}(t)) f_j(\boldsymbol{\omega}(t), \hat{\boldsymbol{\mu}}(t))$

2. Play

$$A_t \leftarrow \begin{cases} \hat{i} = i^*(\hat{\boldsymbol{\mu}}(t)), & \text{if } d\left(\mu_j, \frac{\omega_{\hat{i}} \mu_{\hat{i}} + \omega_j \mu_j}{\omega_{\hat{i}} + \omega_j}\right) > d\left(\mu_{\hat{i}}, \frac{\omega_{\hat{i}} \mu_{\hat{i}} + \omega_j \mu_j}{\omega_{\hat{i}} + \omega_j}\right) \\ C_t, & \text{otherwise.} \end{cases}$$



Frank-Wolfe algorithm (FW)

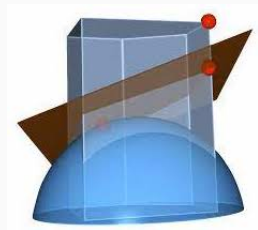
In the view of updating $\omega(t)$, BC corresponds to FW iteration as if the objective function is smooth (**unfortunately, it is not**)

FW for $\max_{\mathbf{x} \in \Sigma} F(\mathbf{x})$ when F is smooth

Take $\mathbf{x}(1) \in \Sigma$ arbitrarily

For $t = 1, \dots, T$ do:

1. $\mathbf{z}(t+1) \leftarrow \operatorname{argmax}_{\mathbf{z} \in \Sigma} \langle \mathbf{z}, \nabla F(\mathbf{x}(t)) \rangle$
2. $\mathbf{x}(t+1) \leftarrow \frac{t}{t+1} \mathbf{x}(t) + \frac{1}{t+1} \mathbf{z}(t+1)$



Curvature

For a compact set \mathcal{K} and a concave function $\psi : \mathcal{K} \mapsto \mathbb{R}$, we define

$$C_\psi(\mathcal{K}) = \sup_{\substack{\mathbf{x}, \mathbf{z} \in \mathcal{K} \\ \alpha \in (0,1] \\ \mathbf{y} = \mathbf{x} + \alpha(\mathbf{z} - \mathbf{x})}} \min_{h \in \partial\psi(\mathbf{x})} \frac{1}{\alpha^2} [\psi(\mathbf{x}) - \psi(\mathbf{y}) + \langle \mathbf{y} - \mathbf{x}, h \rangle] \quad (1)$$

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When \mathcal{K} is a convex domain, a finite curvature permits the convergence of FW (see e.g. [Jag13]). The intuition is that $C_\psi(\mathcal{K})$ provides a controlled bound for each iteration as

$$\psi(\mathbf{x}) + \langle \mathbf{y} - \mathbf{x}, h \rangle - \frac{C_\psi(\mathcal{K})}{\alpha^2} \leq \psi(\mathbf{y}) \leq \psi(\mathbf{x}) + \langle \mathbf{y} - \mathbf{x}, h \rangle,$$

where $h \in \partial\psi(\mathbf{x})$ is the one attaining minimum in (1)



Why does BC fail to reach the optimal allocation?

BC faces three issues:

- (i). F_μ is not smooth
- (ii). Each f_j has an unbounded curvature close to the boundary of Σ
- (iii). μ is unknown initially

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We devise a simple algorithm (FW-based) to track $\mathbf{x}(\mathbf{t}) \xrightarrow{t \rightarrow \infty} \omega^*(\mu)$ by circumventing these issues.



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- Under mild assumption, we show F_μ is the minimum of a finite number of smooth concave functions f_j by *envelop theorem*
- Leveraging this fact, we have a novel and computational efficient construction which continuously approximates the non-smooth points

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Unbounded curvature and unknown μ

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Let updated direction $z(t)$ cover $\mathbf{e}_1, \dots, \mathbf{e}_K$ sufficiently often so that the tracked allocation, $\mathbf{x}(t)$, is kept away from the boundary and each action is forced to be played frequently enough



Assumption 1 and an example

Assumption 1

$\forall i \in \mathcal{I}, \mathcal{S}_i = \{\mu \in \Lambda : i^*(\mu) = i\}$ is open and its complementary $\Lambda \setminus \mathcal{S}_i$ is a finite union of convex set. Namely, a finite collection \mathcal{J}_i of convex set \mathcal{C}_j^i s.t. $\Lambda \setminus \mathcal{S}_i = \cup_{j \in \mathcal{J}_i} \mathcal{C}_j^i$

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Example: BAI for unstructured bandit

Here $\Lambda = \{\boldsymbol{\mu} \in (0, 1)^K : \exists i \in [K] \text{ s.t. } \mu_i > \mu_k, \forall k \neq i\}$, and for each $i \in [K], \mathcal{S}_i = \{\boldsymbol{\mu} \in \Lambda : \mu_i > \mu_k, \forall k \neq i\}$

We can see that $\Lambda \setminus \mathcal{S}_i = \cup_{j \neq i} \mathcal{C}_j^i$, where $\mathcal{C}_j^i = \{\boldsymbol{\lambda} \in \Lambda : \lambda_j > \lambda_i\}$ is a convex set $\forall j \neq i$

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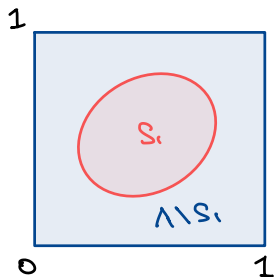
With Assumption 1, we define

$f_j(\boldsymbol{\omega}, \boldsymbol{\mu}) = \inf_{\boldsymbol{\lambda} \in \mathcal{C}_j^i} \sum_{k=1}^K \omega_k d(\mu_k, \lambda_k)$ for any $(\boldsymbol{\omega}, \boldsymbol{\mu}) \in \overset{\circ}{\Sigma} \times \mathcal{S}_i$ and $j \in \mathcal{J}_i$, where $\overset{\circ}{\Sigma}$ is the interior of Σ



A counterexample for Assumption 1

Though most pure exploration and structures satisfy **Assumption 1**, it may not hold for an arbitrary parameter set. For example,



Proposition 1.

Let $i \in \mathcal{I}$, $j \in \mathcal{J}_i$. Define for all $(\omega, \mu) \in \Sigma \times \mathcal{S}_i$,

$$\overline{\lambda_j(\omega, \mu)} = \arg \min_{\lambda \in \text{cl}(\mathcal{C}_j^i)} \sum_k \omega_k d(\mu_k, \lambda_k), \quad (2)$$

where $\text{cl}(\mathcal{C}_j^i)$ is the closure of \mathcal{C}_j^i . Then under Assumption 1, $\overline{\lambda_j(\omega, \mu)}$ is unique for all $(\omega, \mu) \in \mathring{\Sigma} \times \mathcal{S}_i$. In addition, f_j is continuously differentiable on $\mathring{\Sigma} \times \mathcal{S}_i$, and $\forall (\omega, \mu) \in \mathring{\Sigma} \times \mathcal{S}_i$,

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Example (Unstructured BAI)

$$\begin{aligned} f_j(\omega, \mu) &= \omega_j d\left(\mu_j, \frac{\omega_{i^*} \mu_{i^*} + \omega_j \mu_j}{\omega_{i^*} + \omega_j}\right) + \omega_{i^*} d\left(\mu_{i^*}, \frac{\omega_{i^*} \mu_{i^*} + \omega_j \mu_j}{\omega_{i^*} + \omega_j}\right) \\ \nabla_{\omega} f_j(\omega, \mu) &= d\left(\mu_j, \frac{\omega_{i^*} \mu_{i^*} + \omega_j \mu_j}{\omega_{i^*} + \omega_j}\right) \mathbf{e}_j + d\left(\mu_{i^*}, \frac{\omega_{i^*} \mu_{i^*} + \omega_j \mu_j}{\omega_{i^*} + \omega_j}\right) \mathbf{e}_{i^*} \end{aligned}$$

Envelop theorem

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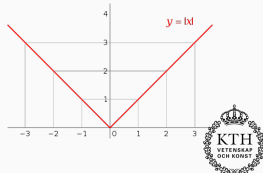
Hence,

- Once we can solve (2), we have $\nabla_{\omega} f_j$ and f_j
- $F_{\mu} = \min_{j \in \mathcal{J}} f_j(\cdot, \mu)$ is the minimum of finite set of smooth

Non-smooth point may cause infinite curvature

$$\text{Recall that } C_\psi(\mathcal{K}) = \sup_{\substack{\mathbf{x}, \mathbf{z} \in \mathcal{K} \\ \alpha \in (0,1] \\ \mathbf{y} = \mathbf{x} + \alpha(\mathbf{z} - \mathbf{x})}} \min_{h \in \partial\psi(\mathbf{x})} \frac{1}{\alpha^2} [\psi(\mathbf{x}) - \psi(\mathbf{y}) + \langle \mathbf{y} - \mathbf{x}, h \rangle]$$

Suppose $\psi(x) = |x|$ over $\mathcal{K} = [-1, 1]$. Let $x \in (0, \frac{1}{2})$, $z = 1 - x$ and $\alpha = 2x$. By definition, $\partial\psi(x) = \{1\}$



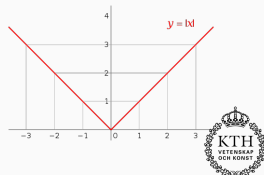
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$$\begin{aligned} C_\psi(\mathcal{K}) &\geq \min_{h \in \partial\psi(x)} \frac{1}{\alpha^2} [\psi(x) - \psi(y) + \langle y - x, h \rangle] \\ &= \frac{1}{4x^2} (|-x| - |x| + 2x) = \frac{1}{2x}. \end{aligned}$$

Hence, $C_\psi(\mathcal{K}) = \infty$



Existing FW works for non-smooth functions

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- enlarging the set of differential [RCS19, CL18] (collect the set of the subdifferentials in the neighborhood)



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Nevertheless, they are all too **computational expensive** to solve in our problem



r -subdifferential space

Recall that $F_{\mu}(\omega) = \min_{j \in \mathcal{J}_i} f_j(\omega, \mu)$ for $\mu \in \mathcal{S}_i$, we define

$$H_{F_{\mu}}(\omega, r) = \text{cov} \{ \nabla_{\omega} f_j(\omega, \mu) : j \in \mathcal{J}_i, f_j(\omega, \mu) < F_{\mu}(\omega) + r \}, \forall r \in \mathbb{R}_+,$$

where $\text{cov}(S)$ is the convex hull of a set S .



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where $\text{cov}(S)$ is the convex hull of a set S .

The advantages include:

- $H_{F_\mu}(\omega, r)$ is a continuous set-mapping
- $H_{F_\mu}(\omega, r)$ is very easy to be computed



r -subdifferential space

Recall that $F_\mu(\omega) = \min_{j \in \mathcal{J}_i} f_j(\omega, \mu)$ for $\mu \in \mathcal{S}_i$, we define

$$H_{F_\mu}(\omega, r) = \text{cov} \{ \nabla_\omega f_j(\omega, \mu) : j \in \mathcal{J}_i, f_j(\omega, \mu) < F_\mu(\omega) + r \}, \forall r \in \mathbb{R}_+,$$

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Our modified FW update is

$$\begin{cases} \mathbf{z}(t+1) \leftarrow \operatorname{argmax}_{\mathbf{z} \in \Sigma} \min_{h \in H_{F_\mu}(\mathbf{x}(t), r_t)} \langle \mathbf{z} - \mathbf{x}(t), h \rangle, \\ \mathbf{x}(t+1) \leftarrow \frac{t}{t+1} \mathbf{x}(t) + \frac{1}{t+1} \mathbf{z}(t+1) \end{cases}$$



Input: Confidence level δ , sequence $\{r_t\}_{t \geq 1}$

Initialization: Sample each arm once and update $\omega(K)$, $\mathbf{x}(K) = (\frac{1}{K}, \dots, \frac{1}{K})$, and $\hat{\mu}(K)$
 $t \leftarrow K$

While $t F_{\hat{\mu}(t)}(\omega(t)) < \beta(\delta, t)$ \leftarrow **Stopping criteria** or $\hat{\mu}(t-1) \notin \Lambda$

IF $\sqrt{[t/K]} \in \mathbb{N}$ or $\hat{\mu}(t-1) \notin \Lambda$, (Forced exploration) $\mathbf{z}(t) \leftarrow (\frac{1}{K}, \dots, \frac{1}{K})$

Else, (FW update)

$$\mathbf{z}(t) \leftarrow \operatorname{argmax}_{\mathbf{z} \in \Sigma} \min_{h \in H_{F_{\hat{\mu}(t-1)}}(\mathbf{x}(t-1), r_t)} \langle \mathbf{z} - \mathbf{x}(t-1), h \rangle$$

$$\text{Update } \mathbf{x}(t) \leftarrow \frac{t-1}{t} \mathbf{x}(t-1) + \frac{1}{t} \mathbf{z}(t)$$

Sample $A_t \leftarrow \operatorname{argmax}_k \mathbf{x}_k(t) / \omega_k(t-1)$ (ties broken arbitrarily)

Update $\omega(t)$ and $\hat{\mu}(t)$

Output: $i^*(\hat{\mu}(t))$



Theoretical Results

Assumption 2

For $\gamma \in (0, \frac{1}{K})$, let $\Sigma_\gamma = \{\omega \in \Sigma : \min_k \omega_k \geq \gamma\}$

Assumption 2.

For all $\mu \in \Lambda$, there exist $L, D > 0$ s.t

- (i). $\forall j \in \mathcal{J}_{i^*}(\mu), \omega \in \Sigma, \|\nabla_\omega f_j(\omega, \mu)\|_\infty \leq L$
- (ii). $\forall \gamma \in (0, 1/K)$ and $\forall j \in \mathcal{J}_{i^*}(\mu), C_{f_j(\cdot, \mu)}(\Sigma_\gamma) \leq \frac{D}{\gamma}$

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We also provide a generic method to verify Assumption 2.



Theorem 1

Consider the FWS algorithm with a sequence $\{r_t\}_{t \geq 1}$ strictly positive reals satisfying

- $\lim_{t \rightarrow \infty} \frac{1}{t} \sum_{s=1}^t r_s = 0,$
- $\lim_{t \rightarrow \infty} t r_t = \infty.$

Under Assumptions 1, 2., the algorithm terminates in finite time a.s. and is δ -PAC. Its sample complexity τ satisfies: $\forall \mu \in \Lambda,$

$$\mathbb{P}_{\mu} \left[\overline{\lim}_{\delta \rightarrow 0} \frac{\tau}{\log(1/\delta)} \leq T^*(\mu) \right] = 1, \text{ and } \overline{\lim}_{\delta \rightarrow 0} \frac{\mathbb{E}_{\mu} [\tau]}{\log(1/\delta)} \leq T^*(\mu).$$

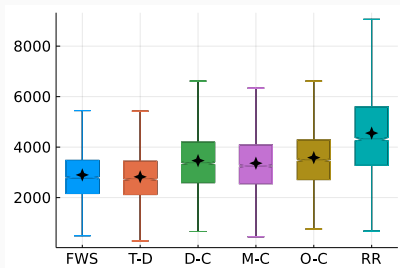
With further assumptions, we can provide **non-asymptotic** upper bound for $\mathbb{E}_{\mu}[\tau]$



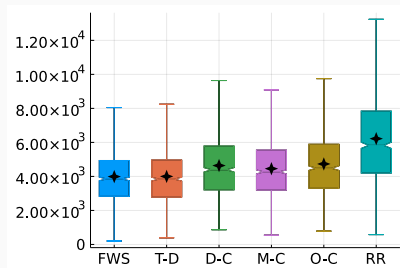
Numerical Results

Experiment (i) Unstructured bandits

Averaged sample complexity at $\delta = 0.01$



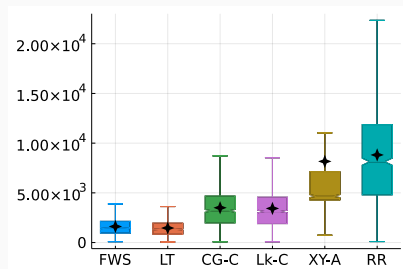
Bernoulli



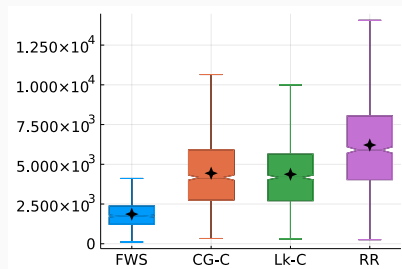
Gaussian

Experiment (ii) Linear bandits

Averaged sample complexity at $\delta = 0.01$



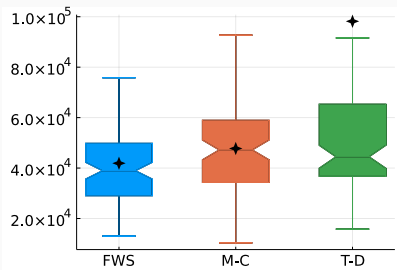
BAI



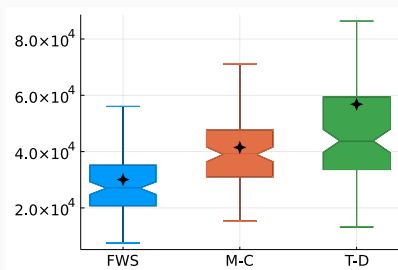
ThresholdingBandit

Experiment (iii) Lipschitz bandits

Averaged Sample complexity at $\delta = 0.01$



Experiment 1



Experiment 2

This is the first result for Lipschitz bandits in literatures

Related work and conclusion

Related works:

- TaS [GK16]: Compute and track the optimal allocation
- LMA [Mén19]: Apply mirror ascent to update $\mathbf{x}(t)$
- Gamification [DMSV20, Sha21, JMKK21]: Use 2 player game to reach $\omega^*(\mu)$

Unclear to extend the above approaches to general structures

Related work and conclusion

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Unclear to extend the above approaches to general structures

Conclusion:

- FWS is computationally and statistically efficient for general pure exploration problems
- Theoretically, FWS matches the instance-specific lower bounds
- Numerically, FWS is competitive to the state-of-art algorithms in structured bandits



Reference

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