### Closing the Computational-Statistical Gap in BAI for Combinatorial Semi-bandits **Ruo-Chun Tzeng**<sup>1</sup>, Po-An Wang<sup>1</sup>, Alexandre Proutiere<sup>1</sup> Chi-Jen Lu<sup>2</sup>



Combinatorial BAI with semi-bandit feedback

Input: K arms  $(\nu_k)_{k \in [K]}$  with mean  $\mu \in \mathbb{R}^K$  and  $\mathcal{X} \subseteq \{0, 1\}^K$ . **Example:** Gaussian reward  $\nu_k = \mathcal{N}(\mu_k, 1), \forall k \in$ 



- Rule: At each round  $t \in \mathbb{N}$ , the learner pulls an action  $\mathbf{x}(t) \in \mathcal{X}$  and observe  $y_k(t) \sim \nu_k$  iff  $x_k(t) = 1$ , and returns her estimated best action  $\hat{i} \in \mathcal{X}$  when she decides to terminate at round  $\tau$ .
- Goal: Design a  $\delta$ -PAC learning strategy s.t. the best action  $i^*(\mu) \in \operatorname{argmax} \langle \mu, \mathbf{x} \rangle$ is identified w.p.  $\geq 1 - \delta$  and  $\mathbb{P}_{\mu}[\tau < \infty] = 1$  while minimizing  $\mathbb{E}_{\mu}[\tau]$ .

**Prior works:** a computational-statistical gap

Any  $\delta$ -PAC algorithm satisfies  $\mathbb{E}_{\mu}[\tau] \geq T^{\star}(\mu) \mathrm{kl}(\delta, 1-\delta)$ , where

$$T^{\star}(\mu)^{-1} = \sup_{\omega \in \Sigma} F_{\mu}(\omega) \text{ with } F_{\mu}(\omega) = \inf_{\lambda \in \mathsf{Alt}(\mu)} \sum_{k=1}^{\kappa} \frac{\omega_k(\mu_k - \mu_k)}{2}$$

Track-and-Stop [6] is statistically optimal but requires to repeatedly solve  $T^{\star}(\hat{\mu}(t-1))^{-1} \Rightarrow$  computationally inefficient.

FWS [8] at the FW-update round has to solve a potentially  $\mathcal{O}(2^{\kappa})$  many convex programs  $\Rightarrow$  computationally inefficient

CombGame [7] is MCP-oracle efficient and statistically optimal  $\Rightarrow$  left open the design of an efficient MCP-oracle





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$$(\lambda_k)^2$$

## Designing efficient MCP based on a structural observation Let $f_{\mathbf{x}}(\boldsymbol{\omega}, \boldsymbol{\mu}) = \inf_{\boldsymbol{\lambda} \in \mathbb{R}: \langle \boldsymbol{i}^{\star}(\boldsymbol{\mu}) - \boldsymbol{x}, \boldsymbol{\lambda} \rangle < 0} \sum_{k=1}^{K} \frac{\omega_{k}(\mu_{k} - \lambda_{k})^{2}}{2} \text{ s.t. } F_{\boldsymbol{\mu}}(\boldsymbol{\omega}) = \min_{\boldsymbol{x} \neq \boldsymbol{i}^{\star}(\boldsymbol{\mu})} f_{\boldsymbol{x}}(\boldsymbol{\omega}, \boldsymbol{\mu}).$ Property of $f_x$ and its Lagrangian dual $g_{\omega,\mu}$ : $f_{\boldsymbol{x}}(\boldsymbol{\omega},\boldsymbol{\mu}) = \max_{\alpha \geq 0} g_{\boldsymbol{\omega},\boldsymbol{\mu}}(\boldsymbol{x},\alpha)$

 $g_{\omega,\mu}(\mathbf{x},\alpha)$  is linear in  $\mathbf{x}$  and concave in  $\alpha$ These properties  $\Rightarrow F_{\mu}(\omega) = \min_{\mathbf{x} \neq \mathbf{i}^{*}(\mu)} \max_{\alpha \geq 0} g_{\omega,\mu}(\mathbf{x}, \alpha)$  as a two-player zero-sum game. We not only want to estimate  $F_{\mu}(\omega)$  but also the *equilibrium* action  $\mathbf{x}_{e}$  s.t.  $F_{\mu}(\omega) = \max_{\alpha \geq 0} g_{\omega,\mu}(\mathbf{x}_{e}, \alpha)$ . This rules out many existing results from applying.  $\succ \mathbf{x}_e$  is required to solve  $\max_{\omega \in \Sigma} F_{\mu}(\omega)$  by the first-order methods  $\triangleright$  Last-iterate convergence [1, 3] are mostly for saddle-point problems

#### Algorithm 1: $(\epsilon, \theta)$ -MCP $(\omega, \mu)$

for  $n = 1, 2, \cdots$  do (Follow-the-Perturbed-Leader)  $\mathcal{Z}_n$  $oldsymbol{x}^{(n)} \in \operatorname*{argmin}_{x 
eq oldsymbol{i}^{\star}(oldsymbol{\mu})} igg( \sum_{m=1}^{n-1} g_{oldsymbol{\omega},oldsymbol{\mu}}$ (Best-Response)  $\alpha^{(n)} \in \operatorname{argmax} g$ if  $\sqrt{n} > \frac{c_{\theta}(1+\epsilon)}{\epsilon \hat{F}}$ , where  $\begin{cases} \hat{F} = \\ n_{\star} \in \end{cases}$ then return  $(\hat{F}, \mathbf{x}^{(n_{\star})})$ ;

end

(Theorem 1) Let 
$$\epsilon, \theta \in (0, 1)$$
 and  $(\omega, \mu) \in \Sigma_+ \times \Lambda$ .  
 $\blacktriangleright \mathbb{P}_{\mu} \Big[ F_{\mu}(\omega) \leq \hat{F} \leq (1 + \epsilon) F_{\mu}(\omega) \Big] \geq 1 - \theta$   
 $\triangleright$  the number of  $i^*$ -oracle calls:  $\mathcal{O} \left( \frac{K^3 D^5 \ln K \ln \theta^{-1} ||\mu||_{\infty}^4 ||\omega^{-1}||_{\infty}^2}{\epsilon^2 F_{\mu}(\omega)^2} \right)$ 





(known by [2])

(our observation)

$$\eta_n \sim \exp(1)^K \text{ and } \eta_n = rac{c_0}{\sqrt{n}}$$
  
 $\sigma_{\mathcal{S},\mu}(\mathbf{x}, \alpha^{(m)}) + rac{\langle \boldsymbol{\mathcal{Z}}_n, \mathbf{x} \rangle}{\eta_n} \end{pmatrix}$ 

$$\mathbf{x}_{\boldsymbol{\omega},\boldsymbol{\mu}}(\mathbf{x}^{(n)},\alpha)$$

$$= g_{\boldsymbol{\omega},\boldsymbol{\mu}}(\boldsymbol{x}^{(n_{\star})}, \alpha^{(n_{\star})})$$
  
= argmin<sub>m \leq n</sub> g\_{\boldsymbol{\omega},\boldsymbol{\mu}}(\boldsymbol{x}^{(m)}, \alpha^{(m)})



#### Our Perturbed Frank-Wolfe Sampling (P-FWS)

We use stochastic smoothing [5, 4] to overcome the nonsmoothness of  $F_{\mu}$  as: all we need is  $i^*$ -oracle and its required graident can be evaluated by envelope theorem [8]. The smoothed objective  $\overline{F}_{\mu,\eta}(\omega) = \mathbb{E}_{\mathbb{Z}\sim \text{Uniform}(B_2)}[F_{\mu}(\omega + \eta \mathbb{Z})]$  satisfies:  $\blacktriangleright \nabla \bar{F}_{\mu,\eta}(\boldsymbol{\omega}) = \mathbb{E}_{\boldsymbol{\mathcal{Z}} \sim \text{Uniform}(B_2)}[\nabla F_{\mu}(\boldsymbol{\omega} + \eta \boldsymbol{\mathcal{Z}})]$  $\blacktriangleright \bar{F}_{\mu,\eta} \text{ is } \frac{\ell K}{n} \text{-smooth and } \bar{F}_{\mu,\eta}(\omega) \xrightarrow{\eta\downarrow 0} F_{\mu}(\omega)$ 

#### High-level design of P-FWS

pull each  $\boldsymbol{x} \in \mathcal{X}_0$  once

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Let  $\mathcal{X}_0$  be a set s.t.  $\forall k \in [K]$ , there exists  $\mathbf{x} \in \mathcal{X}_0$  s.t.  $x_k = 1$ .

P-FWS alternate between two phases:

(to avoid high cost and boundary cases)  $\left( \begin{array}{c} \mathsf{pull} \ \pmb{x}(t) \in \mathsf{argmax}_{\pmb{x} \in \mathcal{X}} \left\langle \nabla \bar{F}_{\hat{\pmb{\mu}}(t-1),\eta_t}(\hat{\pmb{\omega}}(t-1)), \pmb{x} \right\rangle \text{ (ideal FW update)} \right)$ 

(Theorem 2) Let  $\mu \in \Lambda$  and  $\delta \in (0, 1)$ . P-FWS is  $\delta$ -PAC, finishes in finite time,  $\mathbb{P}_{\mu}[\limsup_{\delta \to 0} \frac{\tau}{\ln \delta^{-1}} \leq T^{\star}(\mu)] = 1, \mathbb{E}_{\mu}[\tau] \text{ is bounded by Poly}(K) \text{ in }$ moderate-confidence regime and achieves the minimal in high-confidence regime, and the total number of  $i^*$ -oracle calls is bounded by Poly(K).

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